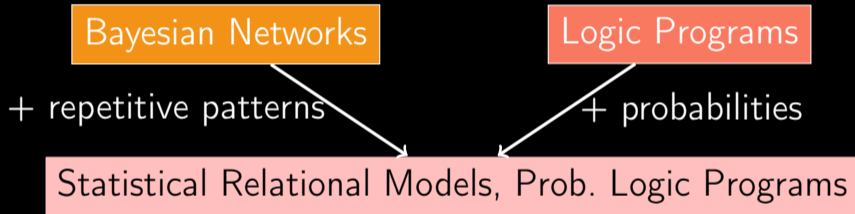


Lecture 2

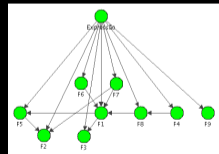
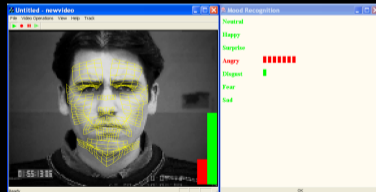
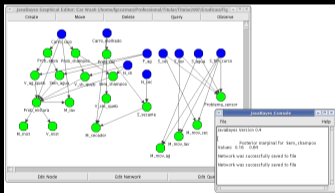
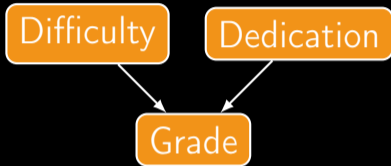
Section 2: Probabilistic Logic Programming

Fabio G. Cozman  
Universidade de São Paulo - Brazil

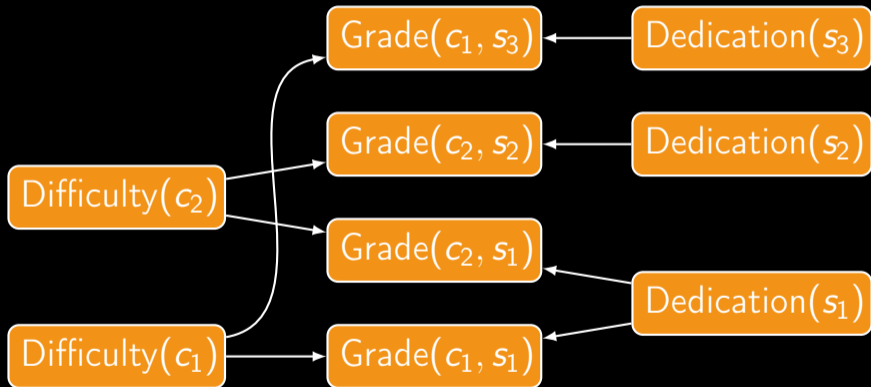
# Mixing things...



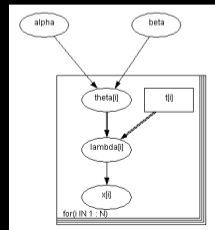
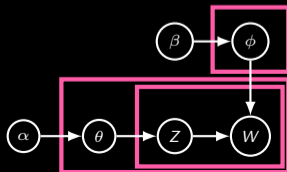
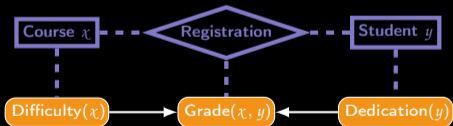
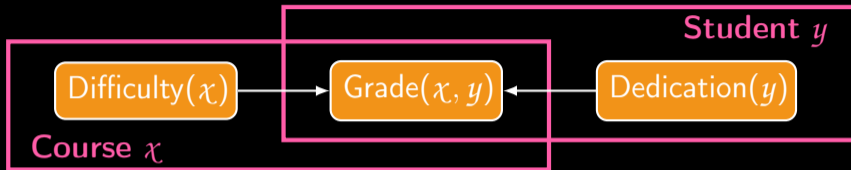
# The propositional language of Bayesian networks



# Bayesian networks: repetitive patterns



# Plates and Probabilistic Relational Models



# Logic programming

- ▶ Rules:

`pass(X, Y) :- student(X), dedicated(X), course(Y), not easy(Y).`

- ▶ Facts:

`edge(a, b).`

Logic programming under uncertainty has been around for a while...

- ▶ Example (Ng and Subrahmanian 1992):

$\text{path}(X, Y): [0.85, 1] \leftarrow \text{connect}(X, Z) \wedge \text{path}(Z, Y): [0.75, 1].$

# The distribution semantics and its variants

- ▶ Dantsin (1990), Poole (1993), Fuhr (1995), Sato (1995):

- ▶ *Probabilistic facts:*

$$\mathbb{P}(A) = 0.2, \mathbb{P}(B) = 0.3.$$

- ▶ A (definite/acyclic/stratified) logic program:

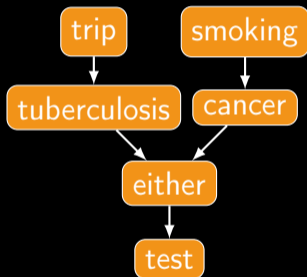
$$C \leftarrow A, B.$$

- ▶ A single probability distribution is induced over programs.



# Acyclic propositional ProbLog (Bayesian network)

0.01 :: trip.  
0.5 :: smoking.  
tuberculosis :- trip, a1.  
tuberculosis :- not trip, a2.  
0.05 :: a1.      0.01 :: a2.  
cancer :- smoking, a3.  
cancer :- not smoking, a4.  
0.1 :: a3.      0.01 :: a4.  
either :- tuberculosis.  
either :- cancer.  
test :- either, a5.      0.98 :: a5.  
test :- either, a6.      0.05 :: a6.



# Acyclic relational ProbLog

0.6 :: dedicated( $X$ ).

0.4 :: hard( $Y$ ).

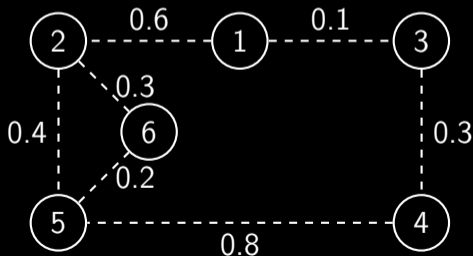
fails( $X$ ,  $Y$ ) :- student( $X$ ), dedicated( $X$ ),  
course( $Y$ ), not hard( $Y$ ), a( $X$ ,  $Y$ ).

0.1 :: a( $X$ ,  $Y$ ).

# Stratified programs

```
edge(X, Y) :- edge(Y, X).  
path(X, Y) :- edge(X, Y).  
path(X, Y) :- edge(X, Z), path(Z, Y)..
```

```
0.6 :: edge(1, 2).  
0.1 :: edge(1, 3).  
0.4 :: edge(2, 5).  
0.3 :: edge(2, 6).  
0.3 :: edge(3, 4).  
0.8 :: edge(4, 5).  
0.2 :: edge(5, 6).
```



# Cyclic programs: many/one/no stable models...

$\text{engineer}(X) :- \text{person}(X), \text{not lawyer}(X).$

$\text{lawyer}(X) :- \text{person}(X), \text{not engineer}(X).$

$0.9 :: \text{person}(\text{dilbert}).$

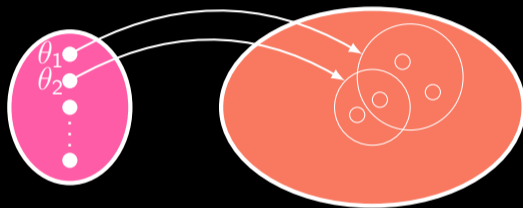
- ▶  $\text{person}(\text{dilbert})$  is false: a unique stable model.
- ▶  $\text{person}(\text{dilbert})$  is true: there are two stable models (one where dilbert is a lawyer, the other in which he is an engineer).

# Credal semantics: intuition

- ▶ One idea: when there are many possible stable models, impose a probability distribution over them (P-log solution: uniform).
  - ▶ For instance: when dilbert is a person, probability 0.5 that it is a lawyer/engineer.
- ▶ Another idea: adopt the set of all possible probability distributions!
  - ▶ T. Lukasiewicz (2005, 2007).

# Credal semantics

- ▶ A total choice may induce a program with many stable models.



- ▶ Probability of each total choice may be distributed freely over answer sets (we have a random set).
- ▶ Semantics is a set of probability distributions that dominates an infinitely-monotone capacity.

# Consequences

- ▶ The credal semantics of a *consistent* PLP is a convex closed set of probability measures (a credal set).
- ▶  $\bar{\mathbb{P}}$  is submodular:  $\bar{\mathbb{P}}(A \vee B) \leq \bar{\mathbb{P}}(A) + \bar{\mathbb{P}}(B) - \bar{\mathbb{P}}(A \wedge B)$ .
- ▶ There are easy formulas for expectations and conditioning:

$$\underline{\mathbb{P}}(A|B) = \frac{\underline{\mathbb{P}}(A \wedge B)}{\underline{\mathbb{P}}(A \wedge B) + \bar{\mathbb{P}}(\neg A \wedge B)}.$$

# Three-coloring

```
red(X) ∨ green(X) ∨ blue(X) :- node(X).  
:- edge(X, Y), red(X), red(Y).  
:- edge(X, Y), green(X), green(Y).  
:- edge(X, Y), blue(X), blue(Y).
```



# Three-coloring for random graphs

$\text{red}(X) \vee \text{green}(X) \vee \text{blue}(X) :- \text{node}(X).$

$:- \text{edge}(X, Y), \text{red}(X), \text{red}(Y).$

$:- \text{edge}(X, Y), \text{green}(X), \text{green}(Y).$

$:- \text{edge}(X, Y), \text{blue}(X), \text{blue}(Y).$

$0.6 :: \text{edge}(1, 2).$

$0.1 :: \text{edge}(1, 3).$

$0.4 :: \text{edge}(2, 5).$

$0.3 :: \text{edge}(2, 6).$

$0.3 :: \text{edge}(3, 4).$

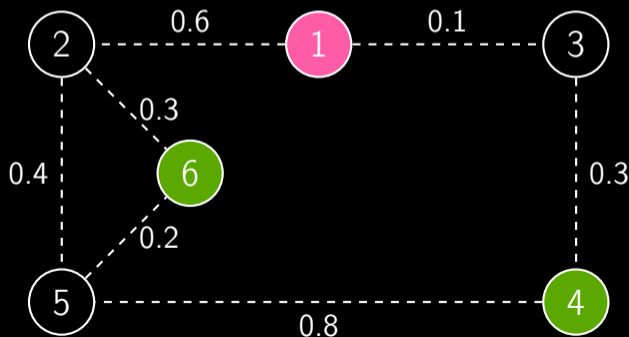
$0.8 :: \text{edge}(4, 5).$

$0.2 :: \text{edge}(5, 6).$

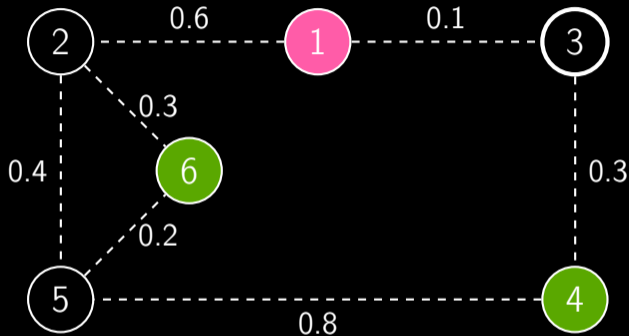
$\text{red}(1).$

$\text{green}(4).$

$\text{green}(6).$



We have:



- ▶  $\overline{\mathbb{P}}(\text{blue}(3)) = 1, \underline{\mathbb{P}}(\text{blue}(3)) = 0.03.$
- ▶  $\overline{\mathbb{P}}(\text{red}(3)) = 0.9, \underline{\mathbb{P}}(\text{red}(3)) = 0.$

# Inference: computing upper probabilities

- ▶ Intuition: for each total choice, we must run brave inference.
- ▶ Algorithm:
  1. Ground the program.
  2. Turn propositional program into CNF formula in propositional logic.
  3. Run counting algorithm for  $\exists$ SAT (Aziz et al. 2015).

## Closing this section...

- ▶ Probabilistic Answer Set Programming gives us a very powerful modeling language.
- ▶ The basic (acyclic non-disjunctive) version offers a “relational” version of Bayesian networks;
- ▶ the general (cyclic, disjunctive) version asks for a semantics based on credal sets.